



IBM Software Group

An Invitation to Ada 2005

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Rational software



@.business on demand software

Agenda

- Object-Oriented Programming
 - Access Types
 - Real-Time, Safety and Criticality
 - General-Purpose Capabilities
 - Predefined Environment and Interfacing



Object-Oriented Programming

- Preserve Ada's strengths for the construction of safe systems
 - ▶ Distinction between specific and class-wide types
 - ▶ Static binding by default, dynamic binding only when necessary
 - ▶ Strong boundary around modules

- Enhance object-oriented features
 - ▶ Multi-package cyclic type structures
 - ▶ Multiple-inheritance type hierarchies
 - ▶ Prefix notation
 - ▶ Accidental overloading or overriding
 - ▶ Extension for protected and task types



Multi-Package Cyclic Type Structures

- Impossible to declare cyclic type structures across library package boundaries
 - ▶ Problem existed in Ada 83
 - ▶ More prominent with the introduction of child units and tagged types
 - ▶ Workarounds result in cumbersome code



Multi-Package Cyclic Type Structures: Example

```
with Employees;  
package Departments is  
  type Department is tagged private;  
  procedure Choose_Manager (D : in out Department;  
                           Manager : in out Employees.Employee);  
private  
  type Emp_Ptr is access all Employees.Employee;  
  type Department is tagged record  
    Manager : Emp_Ptr;  
  end record;  
end Departments;
```

Illegal circularity!

```
with Departments;  
package Employees is  
  type Employee is tagged private;  
  procedure Assign_Employee (E : in out Employee;  
                            D : in out Departments.Department);  
private  
  type Dept_Ptr is access all Departments.Department;  
  type Employee is tagged record  
    Department : Dept_Ptr;  
  end record;  
end Employees;
```



Limited With Clauses

- Gives visibility on a *limited view* of a package
 - ▶ Contains only types and nested packages
 - ▶ Types behave as if they were incomplete
 - ▶ Restrictions on the possible usages of a limited view (no use, no renaming, etc.)
 - ▶ Cycles are permitted among limited with clauses
 - ▶ Imply some kind of “peeking” before compiling a package

- Related change: incomplete tagged types
 - ▶ Can be used as a parameter
 - ▶ Always passed by reference
 - ▶ Support for cycles in object-oriented programming



Limited With Clauses (cont'd)

```
package Departments is
  type Department is tagged; limited view: implicit, visible through limited with
end Departments;
```

```
limited with Departments;
package Employees is
  type Employee is tagged private;
  procedure Assign_Employee (E : in out Employee;
                             D : in out Departments.Department);

private
  type Dept_Ptr is access all Departments.Department;
  type Employee is tagged record
    Department : Dept_Ptr;
  end record;
end Employees;
```

```
with Employees;
package Departments is
  type Department is tagged private;
  procedure Choose_Manager (D : in out Department;
                           Manager : in out Employees.Employee);

private
  type Emp_Ptr is access all Employees.Employee;
  type Department is tagged record
    Manager : Emp_Ptr;
  end record;
end Departments;
```



Multiple-Inheritance Type Hierarchies

- Multiple inheritance too heavy for *Ada 95*
- Java and C# have a lightweight multiple inheritance mechanism: interfaces
 - ▶ Relatively inexpensive at execution time
 - ▶ No conflicts due to inheriting code from multiple parents
- Add interfaces, similar to abstract types but with multiple inheritance
 - ▶ May be used as a secondary parent in type derivations
 - ▶ Have class-wide types
 - ▶ Support for composition of interfaces
 - ▶ No components, no objects
- Related change: null procedures
 - ▶ A procedure declared null need not be overridden



Interfaces: Example

```

❖ type Model is interface;
  type Model_Ref is access all Model'Class;

❖ type Observer is interface;
  procedure Notify
    (O : access Observer;
     M : Model_Ref) is abstract;

❖ type View is interface with Observer;
  type View_Ref is access all View'Class;
  procedure Display
    (V : access View;
     M : Model_Ref) is abstract;

❖ type Controller is interface with Observer;
  type Controller_Ref is
    access all Controller'Class;
  procedure Start
    (C : access Controller;
     M : Model_Ref) is abstract;

  procedure Register
    (M : access Model;
     V : View_Ref) is abstract;
  procedure Register
    (M : access Model;
     C : Controller_Ref) is abstract;

```

■ Model-View-Controller Structure

- ▶ Model: data structure being viewed and manipulated
 - ▶ Observer: waits for a change to a model
 - ▶ View: visual display of a model
 - ▶ Controller: supports input devices for a model
- Note composition of interfaces



Interfaces: Example (cont'd)

```

type Device is tagged private;
procedure Input (D : in out Device);

❖ type Mouse is
    new Device and Controller with private;
procedure Input (D : in out Mouse);

procedure Start (D : access Mouse;
                 M : Model_Ref);
procedure Notify (D : access Mouse;
                  M : Model_Ref);

type Two_Button_Mouse is
    new Mouse with private;
procedure Start
    (D : access Two_Button_Mouse;
     M : Model_Ref);

procedure Register_And_Start
    (D : access Mouse'Class;
     M : Model_Ref);
  
```

- Mouse is a concrete type implementing interface Controller
 - ▶ Only one concrete parent, Device
 - ▶ Any number of interface parents
- Mouse inherits operations from all of its parents
 - ▶ May (but need not) override Input
 - ▶ *Must* override Start and Notify
- Two_Button_Mouse inherits all the operations of Mouse
 - ▶ May (but need not) override some of them
- Register_And_Start is a class-wide operation



Prefix Notation

- A call must identify the package in which an operation is declared
 - ▶ Dispatching operations are often implicitly declared
- Class-wide operations not inherited
 - ▶ Declared in the original package where they appear
- Hard to identify the package where an operation is declared
 - ▶ Difficulty compounded by the fact that the choice between dispatching and class-wide may be an implementation detail
 - ▶ Use clauses are unappealing



Prefix Notation (cont'd)

- Add support for the Object.Operation notation common in other object-oriented languages
 - ▶ Only for tagged types and access designating tagged types
 - ▶ Dispatching operations and class-wide operations declared in the same package as the type are eligible
 - ▶ First parameter of the subprogram must be a controlling parameter
 - ▶ Prefix passed as first parameter

```
M : Model_Ref;
V : View_Ref;
C : Controller_Ref;
D : aliased Mouse;
```

```
...
```

```

■ V.Display (M);           -- equivalent to Display (V, M)
D.Start (M);              -- equivalent to Start (D, M)
D.Input;                  -- equivalent to Input (D)
D.Register_And_Start (M); -- equivalent to
                          -- Register_And_Start (D'Access, M);
```



Accidental Overloading or Overriding

- A typographic error may change overriding into overloading or vice-versa
- Optional syntax to specify that a subprogram is an override or an overload
 - ▶ For compatibility, the absence of a qualifier means “don’t know”

```
type Root_Type is new Ada.Finalization.Controlled with ...;
```

❖ overriding

```
procedure Finalize (Object : in out Root_Type); -- OK.
```

```
type Derived_Type is new Root_Type with ...;
```

❖ overriding

```
procedure Finalise (Object : in out Derived_Type); -- Error here.
```

❖ not overriding

```
procedure Do_Something (Object : in out Derived_Type); -- OK.
```



Extensions for Protected and Task Types

- Type extensions might be useful for protected and task types in addition to record types
- Inheriting of code is complex, notably because of the difficulty to specify how guards and barriers are inherited
- Simpler approach: define interfaces for protected and task types
 - ▶ Includes support for composition of interfaces
 - ▶ A protected or task type may implement any number of interfaces
- Proposal still in a state of flux



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Generalized Use of Anonymous Access Types

- Most OO languages allow free conversion of a reference to a subclass to a reference to its superclass
 - ▶ Ada requires explicit conversions which degrade readability
- Allow anonymous access types in all contexts
 - ▶ Avoids most explicit conversions
 - ▶ Avoids proliferation of access types
 - ▶ Unsure about function returning anonymous access types, yet



Generalized Use of Anonymous Access Types: Example

```
type Animal is tagged ...;
type Horse is new Animal with ...;
type Pig is new Animal with ...;

type Acc_Horse is access all Horse'Class;
type Acc_Pig is access all Pig;

Napoleon, Snowball : Acc_Pig := ...;
Boxer, Clover : Acc_Horse := ...;
Animal_Farm : constant array (Positive range <>) of
    access Animal'Class :=
        (Napoleon, Snowball, Boxer, Clover);

type Noah_S_Arch is
    record
        Stallion, Mare : access Horse;
        Boar, Sow      : access Pig;
    end record;
```



Downward Closures for Access to Subprogram Types

- Access-to-subprogram types subject to accessibility checks
 - ▶ Necessary to prevent dangling references
 - ▶ Requires awkward idioms to deal with nested subprograms

```
type Integrand is access function (X : Float) return Float;
```

```
function Integrate (Fn : Integrand; Lo, Hi : Float) return Float;
```

- Anonymous access-to-subprogram types
 - ▶ Cannot be assigned
 - ▶ Cannot be used to create dangling references

```
❖ function Integrate (Fn : access function (X : Float) return Float;  
Lo, Hi : Float) return Float;
```



Constancy and Null Exclusion

- No access-to-constant parameters or discriminants in Ada 95
- Would be useful for:
 - ▶ Declaring controlling parameters of an operation that doesn't modify the designated object
 - ▶ Providing read-only access via a discriminant
 - ▶ Interfacing with other languages
- Literal null disallowed for anonymous access types
 - ▶ Causes confusion
 - ▶ Problematic when interfacing with a foreign language



Constancy and Null Exclusion (cont'd)

- Define an explicit way to exclude nulls from an access subtype
 - ▶ Make existing anonymous access types include null by default
- Provide a mechanism for declaring constant anonymous access types

```
❖ type Non_Null_Ptr is not null access T;
```

```
-- X guaranteed to not be null.
```

```
procedure Show (X : Non_Null_Ptr);
```

```
-- Pass by reference, but don't allow designated object to be updated;
```

```
-- guarantee Y is non-null.
```


```
❖ procedure Pass_By_Ref (Y : not null access constant Rec);
```

```
-- Any pointer to a graph may be passed to the display routine,  
including null.
```

```
❖ procedure Display (W : access Window;  
                    G : access constant Graph'Class);
```



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Ravenscar Profile for High-Integrity Systems

- *De facto* standard defined by the IRTAW
 - ▶ Intended for use in high-integrity system
 - ▶ Makes it possible to use a reduced, reliable run-time kernel
 - ▶ Many capabilities generally useful for other application domains
- Add new restrictions and pragmas Detect_Blocking and Profile
- Define Ravenscar in terms of predefined restrictions and pragmas
 - ▶ Current users of Ravenscar virtually unaffected
 - ▶ Some application domains only need to abide by some of the restrictions, not the whole profile
 - ▶ Implementers may define new profiles for specific needs



Dynamic Ceiling Priorities

- Tasks have dynamic priorities in Ada 95
- Protected objects only have static ceiling priorities
 - ▶ Unfortunate for some applications
- Add attribute Priority
 - ▶ Prefix is a protected object
 - ▶ Gives the ceiling priority of the object
 - ▶ Attribute is a variable: may be updated, providing dynamic behavior
 - ▶ Completes the language in terms of dynamic priorities



Timing Events

- Some scheduling schemes require to execute code at a particular future time
 - ▶ To asynchronously change the priority of a task
 - ▶ To allow tasks to come off the delay queue at a different priority
- High priority “minder” task needed in Ada 95
 - ▶ Inefficient and inelegant
- Add a mechanism to allow user-defined procedures to be executed at a specified time
 - ▶ Without the need to use a task or a delay statement
- Provided by new predefined unit `Ada.Real_Time.Timing_Events`
 - ▶ Limited private type `Timing_Event` represents an event occurring at some time
 - ▶ Time may be absolute or relative
 - ▶ Protected procedure may be used to handle a timing event



Execution-Time Clocks and Budgeting

- Measuring execution time is fundamental for the safe execution of real-time systems
- Use of aperiodic servers to control allocation is becoming common; requires budget control
- New predefined package `Ada.Real_Time.Execution_Time`
 - ▶ Private type `CPU_Time` represents the CPU time consumed by a task
 - ▶ Handler called when a task has consumed a predetermined amount of CPU
 - ▶ Supports CPU-based scheduling
- New predefined package `Ada.Real_Time.Execution_Time.Group_Budgets`
 - ▶ Private type `Group_Budget` represents a CPU budget for use by a group of tasks
 - ▶ Operations to add or remove a task to a group
 - ▶ Operations to query the remaining budget and change the budget
 - ▶ Handler called when a budget has expired



Scheduling Mechanisms

- Ada 95 only has FIFO scheduling
 - ▶ Other policies may be defined by an implementation, but they are not portable
- Other scheduling techniques are used in practice
 - ▶ Round robin
 - ▶ Earliest deadline first
- Round robin is very common and fits well with the current FIFO
- Earliest deadline first is the preferred scheduling mechanism for soft real-time
 - ▶ Much better CPU usage (40% more before deadlines are missed)
- Add a mechanism to mix scheduling techniques in an application



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Access to Private Units in the Private Part

- Impossible to reference a private unit in the private part of a public package
- Private with clause gives visibility at the beginning of the private part

```
private package Claw.Low_Level_Image_Lists is
    ...
end;
```

```
❖ private with Claw.Low_Level_Image_Lists;
package Claw.Image_List is
    ... -- May not use Low_Level_Image_Lists here.
private
    ... -- May use Low_Level_Image_Lists here.
end;
```



Aggregates for Limited Types

- Limited types prevent copying of values
 - ▶ Have limitations unrelated to copying
 - ▶ Aggregates not available: no full coverage checking
- Allow aggregates for limited types
 - ▶ New syntax to force default initialization of some components

```
private protected type Semaphore is ...;  
type Object is limited  
  record  
    Sem: Semaphore;  
    Size : Natural;  
  end record;  
type Ptr is access Object;
```

```
❖ X : Ptr := new Object'(Sem => <>, Size => 0); -- Coverage checking.
```




Pragma Unsuppress

- Some algorithms may depend on the presence of canonical checks
 - ▶ Interactions with pragma Suppress may lead to bugs
- Pragma Unsuppress revokes the permission granted by Suppress

```
function "*" (Left, Right : Saturation_Type) return Saturation_Type is
  pragma Unsuppress (Overflow_Check);
begin
  return Integer (Left) * Integer (Right);
exception
  when Constraint_Error =>
    if (Left > 0 and Right > 0) or (Left < 0 and Right < 0) then
      return Saturation_Type'Last;
    else
      return Saturation_Type'First;
    end if;
end "*";
```



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Unchecked Unions: Variant Records with no Run-Time Discriminant

- No support in Ada 95 for interfacing with C unions
 - ▶ Unchecked_Conversion not satisfactory
- Pragma Unchecked_Union prevents discriminants from being stored
 - ▶ Operations that need to read a discriminant are either illegal or raise Program_Error

```
union {  
  spvalue double;  
  struct {  
    length int;  
    first *double;  
  } mpvalue;  
} number;
```

```
type Number (Kind : Precision) is  
  record  
    case Kind is  
      when Single_Precision =>  
        SPValue : Long_Float;  
      when Multiple_Precision =>  
        MP_Value_Length : Integer;  
        MP_Value_First : Access_Long_Float;  
    end case;  
  end record;  
pragma Unchecked_Union (Number);
```



Vector and Matrix Operations

- ISO/IEC 13813 defined real and complex vectors and matrices for Ada 83
 - ▶ No support for basic linear algebra
 - ▶ Not provided by vendors

- Integrate this capability in Annex G (Numerics)
 - ▶ Two new predefined units: `Ada.Numerics.Generic_Real_Arrays` and `Ada.Numerics.Generic_Complex_Arrays`
 - ▶ Adapted for Ada 05
 - ▶ Add support for basic linear algebra: inversion, resolution, eigensystem
 - ▶ May be used as an interface to existing linear algebra libraries or as a self-standing implementation



Container Library

- Numerous container libraries available as public domain software
 - ▶ Stacks, lists, maps, sets, trees, queues, graphs, etc...
- Language-defined containers would improve portability and usability of the language
- Delegated by the ARG to outside groups
 - ▶ Not enough resources in the ARG to fully specify a bulky API
 - ▶ Proposals will be evaluated by the ARG



Directory Operations

- Modern operating systems have a tree-structured file system
- Applications need to manage these file systems
- New predefined package `Ada.Directory_Operations`
 - ▶ Query and set the current directory
 - ▶ Create and remove directories or directory trees
 - ▶ Copy and rename files and directories
 - ▶ Decompose and compose file and directory paths
 - ▶ Check the existence, size and modification time of a file
 - ▶ Iterate over files and directories



Conclusion

- Snapshot of work in progress
 - ▶ Other features are being considered
 - ▶ More work needed to integrate all the changes together: consistency, orthogonality

- Schedule-driven: expect completion around the end of 2005
 - ▶ Features frozen by the end of this year
 - ▶ Implementers may want to do pilot implementation of some new features, based on user demand

- Make Ada safer, more powerful, more appealing to new and existing users

